

Functional coefficients in solutions of non-commutative differential equations

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Outline

1 Motivation

2 Main Theorem

3 Examples

- Polylogarithms
- Counterexample
- Hyperlogarithms

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- Generalization (in view of multiplications) : **Polyzetas**

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- Polyzetas are values of **polylogs** at 1 :

$$\zeta(\mathbf{s}) = \text{Li}_{\mathbf{s}}(1).$$

Polylogs can be manipulated as **shuffles** : algebra structure.

- (Non commutative) Differential equation :

Drinfel'd equation $\frac{d}{dz} T(z) = \left(\frac{1}{z} x_0 + \frac{1}{1-z} x_1 \right) T(z).$

- Bailey & Borwein & Girgensohn 1994,
- Zagier 1994,
- Flajolet & Salvy 1998,
- Minh & Petitot & Van Der Hoeven 1998,
- Waldshmidt 2000,
- Cartier 2002,
- Racinet 2002,
- Ecalle 2003,
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Examples of relations between **Multiple Zeta Values** (M. Bigotte) :

$$\zeta(7, 5) = -\frac{6117808}{2627625}\zeta(2)^6 - 7\zeta(10, 2) + 28\zeta(5)\zeta(7) + 14\zeta(3)\zeta(9)$$

$$\zeta(9, 3) = -\frac{93976}{79625}\zeta(2)^6 - \frac{9}{2}\zeta(10, 2) + 12\zeta(5)\zeta(7) + 9\zeta(3)\zeta(9)$$

$$\zeta(9, 3) = \frac{9}{14}\zeta(7, 5) + \frac{388112}{1226225}\zeta(2)^6 - 6\zeta(5)\zeta(7)$$

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Notations

If X is an alphabet and k a field,

- X^* denotes the set of words with letters in X ;
- $k\langle X \rangle$ is the algebra of non commutative polynomials with coefficients in k ,
- $k\langle\langle X \rangle\rangle$ the algebra of non commutative series with coefficients in k ;
- If S is a (non commutative) series and w a word, $\langle S|w \rangle$ denotes the coefficient of S on w .

Data

- X an alphabet.
- (\mathbb{A}, d) a commutative differential algebra over the ring k :
 - differential : $\forall a, b \in \mathbb{A}, d(ab) = d(a)b + ad(b)$;
 - d is **linear** over k ;

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- We require that $\ker(d) = k$ (set of constants = k).
- Extension of d to $\mathbb{A}\langle\langle X \rangle\rangle$:

$$\forall S \in \mathbb{A}\langle\langle X \rangle\rangle, d(S) = \sum_{w \in X^*} d(\langle S|w \rangle)w.$$

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- Let \mathbb{D} be a differential subfield of \mathbb{A} (i.e. $d(\mathbb{D}) \subset \mathbb{D}$).
- M : a homogeneous series **of degree 1** :

$$M = \sum_{x \in X} u_x x \in \mathbb{D}_{=1}\langle\langle X \rangle\rangle.$$

Theorem 1 : Linear independence

Suppose that $T \in \mathbb{A}\langle\langle X \rangle\rangle$ is a solution of the differential equation

$$dT = MT; \langle T | 1_{X^*} \rangle = 1.$$

The following conditions are equivalent¹ :

- i) The family $(\langle T | w \rangle)_{w \in X^*}$ of coefficients of T is free over \mathbb{D} .
- ii) The family of coefficients $(\langle T | y \rangle)_{y \in X \cup \{1_{X^*}\}}$ is free over \mathbb{D} .
- iii) The family $(u_x)_{x \in X}$ is such that, for $f \in \mathbb{D}$ and $\alpha_x \in k$

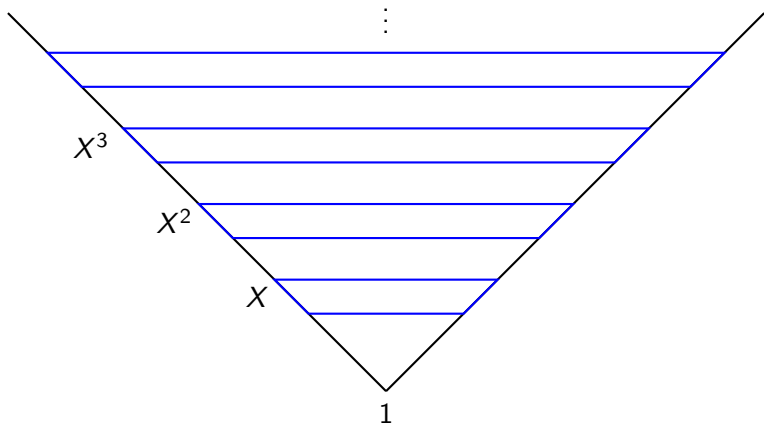
$$d(f) = \sum_{x \in X} \alpha_x u_x \implies (\forall x \in X)(\alpha_x = 0).$$

- iv) The family $(u_x)_{x \in X}$ is free over k and

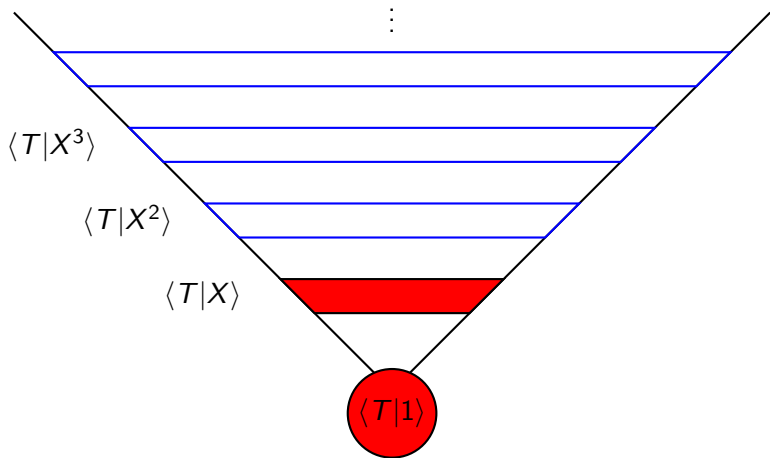
$$d(\mathbb{D}) \cap \text{span}_k((u_x)_{x \in X}) = \{0\}.$$

¹Independence of hyperlogarithms over function fields via algebraic combinatorics, M. D., G. H. E.

“Slices” of the free monoid



Linear independence of the bottom triangle

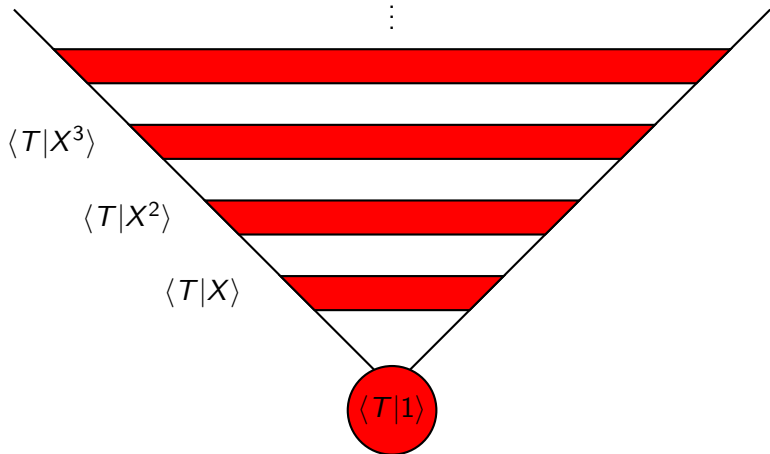


Linear independence of the bottom triangle

\Rightarrow

Linear independence of the whole triangle

\vdots

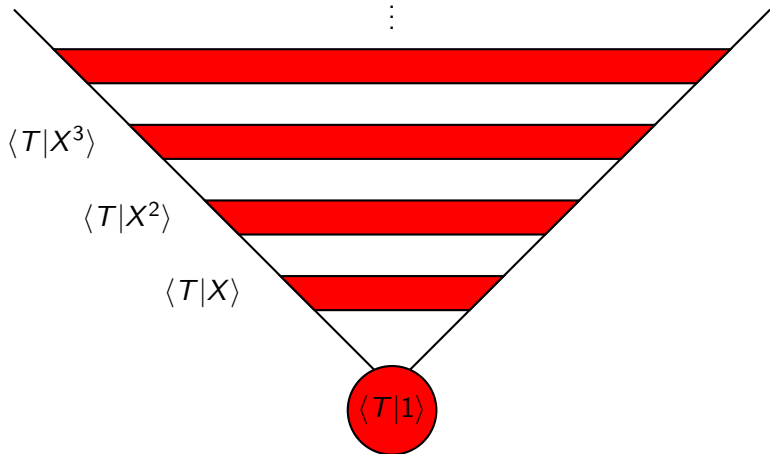


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\Rightarrow Test of linear relations over a **finite number** of functions.

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Polylogarithms

$$X = \{x_0, x_1\}. \quad \Omega = \mathbb{C} \setminus (]-\infty, 0[\cup]1, +\infty[). \quad u_0(z) = \frac{1}{z}, \quad u_1(z) = \frac{1}{1-z}.$$



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Definition

$\forall z \in \Omega,$

$$\text{Li}_{x_0^n}(z) = \frac{\ln^n(z)}{n!}.$$

$$\text{Li}_{x_1 w}(z) = \int_0^z \frac{dt}{1-t} \text{Li}_w(t),$$

and, $\forall w \in X^* x_1 X^*,$

$$\text{Li}_{x_0 w}(z) = \int_0^z \frac{dt}{t} \text{Li}_w(t).$$

Remark

Let $w = x_0^{s_1-1} x_1 \dots x_0^{s_k-1} x_1 \leftrightarrow \mathbf{s} = (s_1, \dots, s_k)$.

It can be shown that the Taylor expansion of these functions is given by

$$\text{Li}_w(z) = \text{Li}_{\mathbf{s}}(z) = \sum_{n_1 > n_2 > \dots > n_k > 0} \frac{z^{n_1}}{n_1^{s_1} \dots n_k^{s_k}}.$$

Application of the theorem

- $\mathbb{A} =$ functions from $\Omega = \mathbb{C} \setminus (]-\infty, 0[\cup]1, +\infty[)$ to \mathbb{C} .
- $\mathbb{D} =$ field of functions on Ω (germs of analytic functions).
- $T =$ generating series of polylogs : $T(z) = \sum_{w \in X^*} \text{Li}_w(z) w$.
- $M(z) = \frac{1}{z} x_0 + \frac{1}{1-z} x_1$.
- Differential equation : **Drinfel'd equation**

$$\frac{d}{dz} T(z) = M(z) T(z).$$

Consequence : Linear independance of polylogs over \mathbb{D} .

Let $X = \{x_0, x_1\}$. $u_0(z) = 1$ and $u_1(z) = \frac{1}{z}$.

Encoding integrals : $x_i \rightarrow \int_{z_i}^z \cdot u_i(s) ds$ $x_i = \alpha_{z_i}^z(x_i)x_i$, $z_0 = 0, z_1 = 1$.

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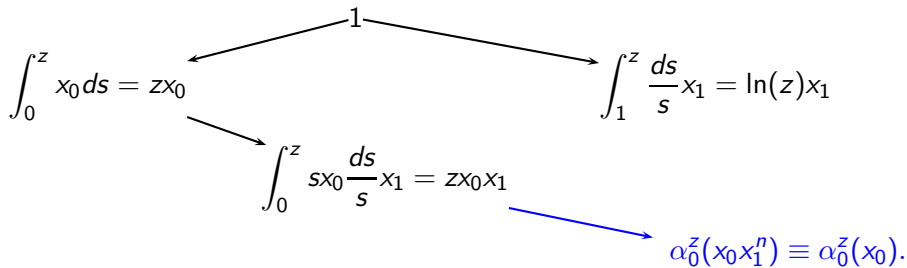
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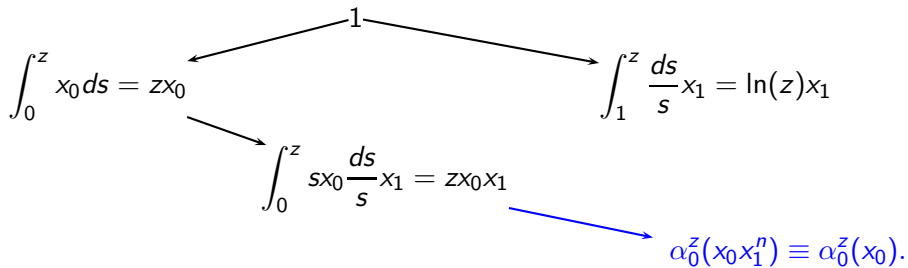
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Problem : Third condition of the theorem.

iii) The family $(u_x)_{x \in X}$ is such that, for $f \in \mathbb{D}$ and $\alpha_x \in k$

$$d(f) = \sum_{x \in X} \alpha_x u_x \implies (\forall x \in X)(\alpha_x = 0).$$

Hyperlogarithms

Definition (1928, Lappo-Danilevski)

Let $a_0, \dots, a_k \in \mathbb{C}$. Then

$$L(a_{i_n}, \dots, a_{i_1} | \gamma) = \int_{z_0}^z \int_{z_0}^{s_n} \dots \int_{z_0}^{s_2} \frac{ds_1}{s_1 - a_{i_1}} \dots \frac{ds_n}{s_n - a_{i_n}}$$

with $\gamma : z_0 \rightsquigarrow z$ a path such that

$$a_{j_i} \notin \gamma \text{ and } s_i \in \gamma, \forall i \in \{1, \dots, n\}.$$

- If $z_0 \neq a_{i_1}$, the integral converges.

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- If $z_0 \neq a_{i_1}$, the integral converges.
- Our theorem applies as well to families of inputs of the type

$$u_i(z) = \frac{\lambda_i}{z - a_i}, \lambda_i \in \mathbb{C}^*.$$

Field of germs

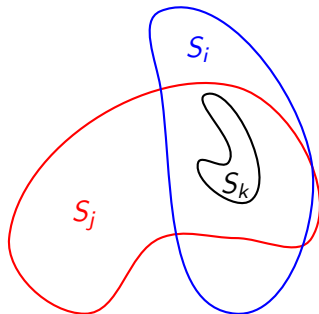
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- Ω : connected, simply connected, analytic domain.
- \mathfrak{B} : a filter basis of Ω of open connected (non void) subsets of Ω :

$$\forall S_i, S_j \in \mathfrak{B}, \exists S_k \in \mathfrak{B}, S_k \subset S_i \cap S_j.$$



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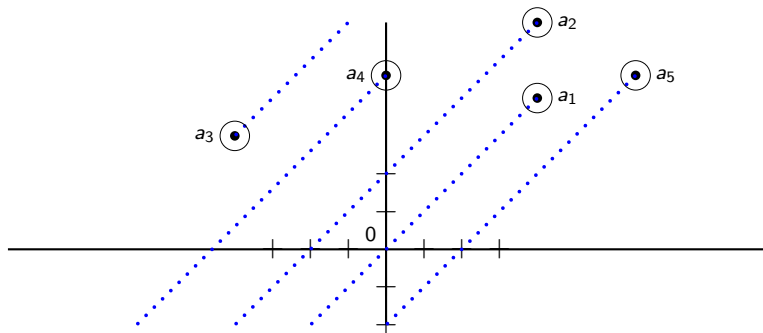
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- **Compatibility with restrictions** : $\forall U, W \in \mathfrak{B}, W \subset U$,

$$\text{res}_{WU}(C[U]) \subset C[W].$$

Example

What is the good domain if the inputs are $u_i(z) = \frac{\lambda_i}{z - a_i}$?

It is always possible to cut the complex plane with half rays to form a **simply connected domain** on which the u_i 's are analytic :



Encoding integrals with words

$$X = \{x_1, \dots, x_n\}.$$

Definition of the **iterated integrals** $\alpha_{z_0}^z(w)$ for $w \in X^*$ and $z_0, z \in \Omega$:

$$\alpha_{z_0}^z(1_{X^*}) = 1;$$

$$\alpha_{z_0}^z(x_i) = \int_{z_0}^z u_i(s) ds, \quad x_i \in X;$$

$$\alpha_{z_0}^z(x_i w) = \int_{z_0}^z u_i(s) ds \alpha_{z_0}^s(w), \quad x_i \in X, w \in X^*.$$

Generating series of hyperlogarithms :

$$T(z) := \sum_{w \in X^*} \alpha_{z_0}^z(w) w.$$

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General idea : Derivating T term by term, we obtain the following **non commutative differential equation** :

$$\frac{d}{dz} T(z) = M(z) T(z), \quad M(z) = \sum_{x_i \in X} u_i(z) x_i.$$

Integrator

Let $M(z) = \sum_{x_i \in X} u_i(z)x_i = \sum_{x_i \in X} \frac{1}{z - a_i} x_i$ and $z_0 \in \mathbb{C}$.

We define the **integrator** H_{z_0} :

$$H_{z_0} : \mathbb{A}\langle\langle X \rangle\rangle \rightarrow \mathbb{A}_{\geq 1}\langle\langle X \rangle\rangle$$

$$S \mapsto H_{z_0}[S] = \int_{z_0}^z M(s)S(s)ds$$

Since $\forall S, H_{z_0}^n[S] \in \mathbb{A}_{\geq n}\langle\langle X \rangle\rangle$,

$$\langle H_{z_0}^n[S] | w \rangle \neq 0 \text{ only for } n \leq |w|.$$

Therefore, we can define the sum

$$\sum_{w \in X^*} \sum_{n \geq 0} \langle H_{z_0}^n[S] | w \rangle w = H_{z_0}^*[S] = \sum_{n \geq 0} H_{z_0}^n[S].$$

(Non commutative) Differential equation

It is clear that

$$H_{z_0}^* = 1 + H_{z_0} H_{z_0}^*$$

Therefore, $\forall S \in \mathbb{A}\langle\langle X \rangle\rangle$ such that $dS = 0$ (constant series),

$$d(H_{z_0}^*[S]) = d(S + H_{z_0}(H_{z_0}^*[S])) = MH_{z_0}^*[S],$$

and $H_{z_0}^*[S]$ satisfies the (non commutative) differential equation

$$dP = MP.$$

Since

$$H_{z_0}^*[1] = T(z) = \sum_{w \in X^*} \alpha_{z_0}^z(w)w$$

we obtain the promised differential equation.

Application of the theorem

- $X = \{x_1, \dots, x_n\}$.
- \mathfrak{B} : filter basis obtained by cutting the complex plane.
- \mathbb{D} = Field of germs of functions on \mathfrak{B} fulfilling condition $i)$ of theorem 1 (e.g., the field of rational functions or the field of functions that are inessential at all the points a_i).
- T = generating series of hyperlogs :
$$T(z) = \sum_{w \in X^*} \alpha_{z_0}^z(w)w.$$
- $$M(z) = \sum_{i=1}^n \frac{1}{z - a_i} x_i, \quad \lambda_i \neq 0, \quad \forall i.$$

Consequence : Linear independance of hyperlogs over \mathbb{D} .

Conclusion and perspectives

Conclusion :

- **New proof** of known results (without monodromy²) ;
- **Generalization** of these results to wider classes of algebras and functions.

Perspectives : Implementation.

²Computation of the monodromy of generalized polylogarithms, Minh & Petitot & Van Der Hoeven, ISSAC'98

Thank you for your attention!