

# Independence of hyperlogarithms over function fields via algebraic combinatorics

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# Outline

## 1 Introduction

- Definition
- Known results
- Encoding

## 2 Differential equation

## 3 Theorems and application to hyperlogarithms

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## Definition (1928, Lappo-Danilevski)

Let  $a_0, \dots, a_k \in \mathbb{C}$ . Then

$$L(a_{i_n}, \dots, a_{i_0} | z, z_0) = \int_{z_0}^z \int_{z_0}^{s_n} \dots \int_{z_0}^{s_1} \frac{ds_0}{s_0 - a_{i_0}} \dots \frac{ds_n}{s_n - a_{i_n}}.$$

- Of order  $n$  ;
- If  $z_0 \neq a_{i_0}$ , the integral converges.

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- If  $z_0 \neq a_{i_0}$ , the integral converges.
- Largest field on which these functions are independent ?
- If no linear independence, describe the relations.

# Polylogarithms

$$X = \{x_0, x_1\}.$$

## Definition

$$\text{Li}_{x_1}(z) = -\ln(1-z) \text{ and } \text{Li}_{x_0}(z) = \ln(z).$$

$$\text{Li}_{x_0 w}(z) = \int_0^z \frac{dt}{t} \text{Li}_w(t); \quad \text{Li}_{x_1 w}(z) = \int_0^z \frac{dt}{1-t} \text{Li}_w(t).$$

$$w = x_0^{s_1-1} x_1 \dots x_0^{s_k-1} x_1 \leftrightarrow \mathbf{s} = (s_1, \dots, s_k).$$

$$\text{Li}_w(z) = \text{Li}_{\mathbf{s}}(z) = \sum_{n_1 > n_2 > \dots > n_k > 0} \frac{z^{n_1}}{n_1^{s_1} \dots n_k^{s_k}}.$$

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$\mathbb{C}$ -linear independence [[Van der Hoeven, Hoang Ngoc Minh, Petitot, 2000](#)].

# Coloured Polylogarithms

## Definition

$$\mathrm{Li}_{\mathbf{b},\mathbf{s}}(z) = \sum_{n_1 > n_2 > \dots > n_k > 0} \frac{b_1^{n_1} \dots b_k^{s_k}}{n_1^{s_1} \dots n_k^{s_k}} z^{n_1}$$

with  $\mathbf{s} = (s_1, \dots, s_k)$ ,  $\mathbf{b} = (b_1, \dots, b_k)$ .

Linear independence over  $\mathbb{C} \left[ z, \frac{1}{z}, \left( \frac{1}{\rho_i^{-1} - z} \right)_{i=0, \dots, n} \right]$  where the  $(\rho_i)_i$ 's are the  $n^{\mathrm{th}}$  roots of 1 [Hoang Ngoc Minh, 2004].

# Encoding integrals with words

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Recursive definition of the **iterated integrals**  $\alpha_{z_0}^z(w)$  for  $w \in X^*$  :

$$\alpha_{z_0}^z(1_{X^*}) = 1;$$

$$\alpha_{z_0}^z(x_i) = \int_{z_0}^z u_i(s) ds, \quad x_i \in X;$$

$$\alpha_{z_0}^z(x_i w) = \int_{z_0}^z u_i(s) ds \alpha_{z_0}^s(w), \quad x_i \in X, w \in X^*.$$

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Then if  $u_i(z) = \frac{1}{z - a_i} \leftrightarrow x_i$ ,

$$\alpha_{z_0}^z(x_{j_0} \dots x_{j_n}) = L(a_{j_n}, \dots, a_{j_0} | z, z_0).$$

**Generating series of hyperlogarithms :**

$$L(z) := \sum_{w \in X^*} \alpha_{z_0}^z(w) w.$$

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**General idea** : Derivating term by term  $L$ , we obtain the following non commutative differential equation :

$$\frac{d}{dz} L(z) = M(z) L(z), \quad \text{with } M(z) = \sum_{x_i \in X} u_i(z) x_i.$$

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# Non commutative polynomials and series

Let  $R$  be a (commutative) **ring** (with unit).

- $R\langle X \rangle$  : non commutative polynomials ;
- $R\langle\langle X \rangle\rangle$  : non commutative series ;
- For  $S \in R\langle\langle X \rangle\rangle$ ,  $\langle S|w \rangle$  is the coefficient of  $S$  on the word  $w$  :

$$S = \sum_{w \in X^*} \langle S|w \rangle w.$$

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Special case :  $R = (\mathfrak{A}, d)$ , a **commutative differential algebra** over  $k$ .

- differential :  $\forall a, b \in \mathfrak{A}, d(ab) = d(a)b + ad(b)$  ;
- $d(a + b) = d(a) + d(b)$  ;
- we require (for the following theorems) that  $\ker(d) = k$  (the only constants for the derivation are the elements of the ground field  $k$ ).

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- $d(a + b) = d(a) + d(b)$  ;
- $\ker(d) = k$ .

We can extend in a unique way  $d$  to  $\mathfrak{A}\langle\langle X \rangle\rangle$  :

$$\forall S \in \mathfrak{A}\langle\langle X \rangle\rangle, d(S) = \sum_{w \in X^*} d(\langle S|w \rangle)w.$$

# Integrator

Let  $M \in \mathfrak{A}_{=1}(\langle X \rangle)$  ( $M(z) = \sum_{x \in X} u_x(z)x$ ) and  $z_0 \in \mathbb{C}$ .

We define the **integrator**  $H_{z_0}$  :

$$H_{z_0} : \mathfrak{A}(\langle X \rangle) \rightarrow \mathfrak{A}_{\geq 1}(\langle X \rangle)$$

$$S \mapsto H_{z_0}[S] = \int_{z_0}^z M(s)S(s)ds$$

with

$$\int_{z_0}^z T(s)ds = \sum_{w \in X^*} \left( \int_{z_0}^z \langle T|w \rangle ds \right) w, \quad \forall T \in \mathfrak{A}(\langle X \rangle).$$

# Iterated integrals

Since  $\forall S, H_{z_0}^n [S] \in \mathfrak{A}_{\geq n}(\langle X \rangle)$ ,

$$\langle H_{z_0}^n [S] | w \rangle \neq 0 \text{ only for } n \leq |w|.$$

Therefore, we can define the sum

$$\sum_{w \in X^*} \sum_{n \geq 0} \langle H_{z_0}^n [S] | w \rangle w = H_{z_0}^* [S] = \sum_{n \geq 0} H_{z_0}^n [S].$$

**Link with hyperlogarithms :** With  $u_i(z) = \frac{1}{z - a_i}, \forall i$ ,

$$H_{z_0}^* [1] = L(z) = \sum_{w \in X^*} \alpha_{z_0}^z(w) w$$

and

$$\langle H_{z_0}^* [1] | w \rangle = \alpha_{z_0}^z(w).$$

## (Non commutative) Differential equation

It is clear that

$$H_{z_0}^* = 1 + H_{z_0} H_{z_0}^*$$

Therefore,  $\forall S \in \mathfrak{A}(\langle X \rangle)$  such that  $dS = 0$  (constant series),

$$d(H_{z_0}^*[S]) = d(S + H_{z_0}(H_{z_0}^*[S])) = MH_{z_0}^*[S],$$

and  $H_{z_0}^*[S]$  satisfies the (non commutative) differential equation

$$dT = MT.$$

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# Theorem 1 : Linear independence

Let  $\mathfrak{C}$  be a differential subfield of  $\mathfrak{A}$  (i.e.  $d(\mathfrak{C}) \subset \mathfrak{C}$ ).

We suppose that  $T \in \mathfrak{A}\langle\langle X \rangle\rangle$  is a solution of the differential equation

$$dT = MT; \langle T | 1_{X^*} \rangle = 1$$

where  $M$  is a homogeneous series of degree 1 :  $M = \sum_{x \in X} u_x x \in \mathfrak{C}_{\geq 1}\langle\langle X \rangle\rangle$ .

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The following conditions are equivalent :

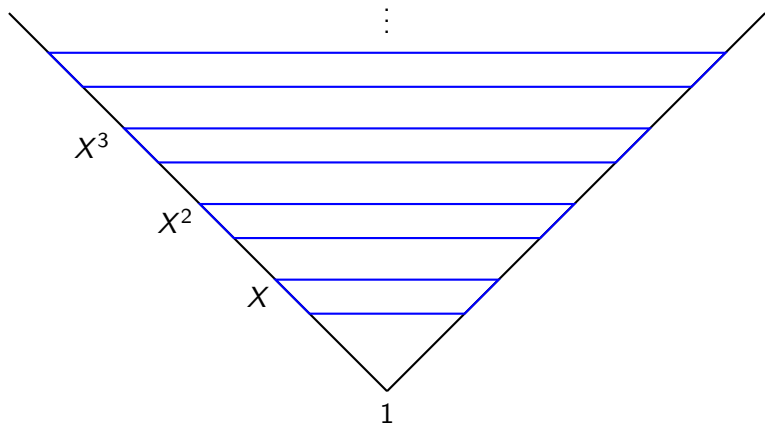
- i) The family  $(\langle T|w \rangle)_{w \in X^*}$  of coefficients of  $T$  is free over  $\mathfrak{C}$ .
- ii) The family of coefficients  $(\langle T|y \rangle)_{y \in X \cup \{1_{X^*}\}}$  is free over  $\mathfrak{C}$ .
- iii) The family  $(u_x)_{x \in X}$  is such that, for  $f \in \mathfrak{C}$  and  $\alpha_x \in k$

$$d(f) = \sum_{x \in X} \alpha_x u_x \implies (\forall x \in X)(\alpha_x = 0).$$

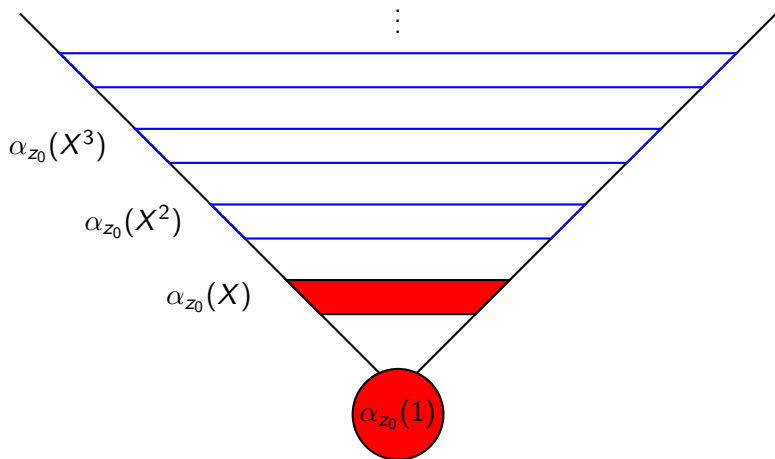
- iv) The family  $(u_x)_{x \in X}$  is free over  $k$  and

$$d(\mathfrak{C}) \cap \text{span}_k((u_x)_{x \in X}) = \{0\}.$$

## “Slices“ of the free monoid



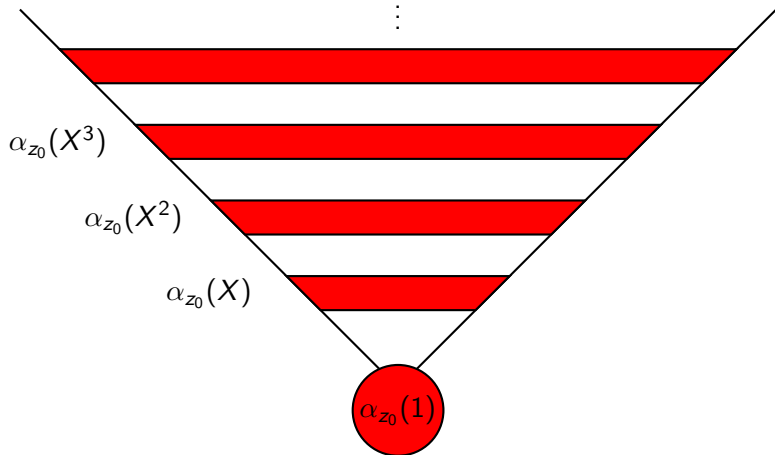
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 $\Rightarrow$ 

Linear independence of the whole triangle

 $\vdots$ 

# Field of germs

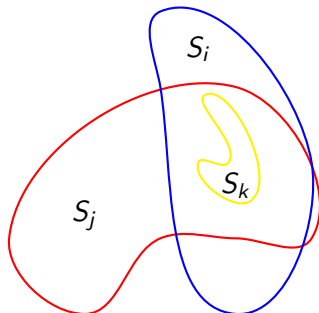
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- $V$  : connected, simply connected, analytic variety :  $\mathbb{C} \setminus \{a_1, \dots, a_n\}$ ;
- $\mathfrak{B}$  : a filter basis of  $V$  of open connected subsets of  $V$  :

$$(\mathbb{C} \setminus (a_j)_{j \in \mathcal{J}})_{\mathcal{J} \in \mathfrak{B}}$$



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 $\forall U \in \mathfrak{B}, C[U]$  is a subring of  $\mathcal{C}^\omega(U, \mathbb{C})$ , satisfying :

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- **Inverse** : if  $f \in C[U] \setminus \{0\}$ ,  $\exists W \in \mathfrak{B}$  such that  $W \subset U - \mathcal{O}_f$  and  $f^{-1}$  (defined on  $W$ ) is in  $C[W]$ .

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- **Compatibility with restrictions** :  $\forall U, W \in \mathfrak{B}, W \subset U$ ,

$$\text{rest}_{WU}(C[U]) \subset C[W].$$

## Theorem 2 : Application to hyperlogarithms

Let

- $M = \sum_{i=1}^n \frac{\lambda_i}{z - a_i} x_i$  be a multiplier on  $X$  with all  $\lambda_i \neq 0$ ,
- $S$  be any regular solution of

$$\frac{d}{dz} S(z) = M(z)S(z).$$

Then, let  $C$  be a differential field of functions defined on  $V$  which **does not** contain linear combinations of logarithms on any domain but which contains  $z$  and the constants (as, for example, the rational functions).

If  $U \in \mathfrak{B}$  and  $P \in C[U] \langle X \rangle$ , one has

$$\langle S | P \rangle = 0 \implies P = 0.$$

**Linear independence on  $C[U]$ .**

# Application : Hyperlogarithms

## Corollary of the previous theorem :

$R$  denotes the ring of functions that can be analytically extended to some  $V \cup U_{a_1} \cup U_{a_2} \cup \dots \cup U_{a_n}$  ( $U_{a_i}$  open neighborhood of  $a_i$ ) and have non-essential singularities at these points.

Then the set of hyperlogarithms  $(\langle S|w \rangle)_{w \in X^*}$  is a set of linearly independent functions over  $R$ .

# Conclusion and perspectives

## Conclusion :

- **New and simpler proof** of known results (without monodromy) ;
- **Generalization** of these results to a wider class of algebras.

## Perspectives :

- Find the maximal field on which the linear independence is preserved ;
- **Write a program** that
  - tests the linear (in)dependence of functions over a given differential ring ;
  - gives the relations between the coefficients if necessary.

Thank you for your attention!